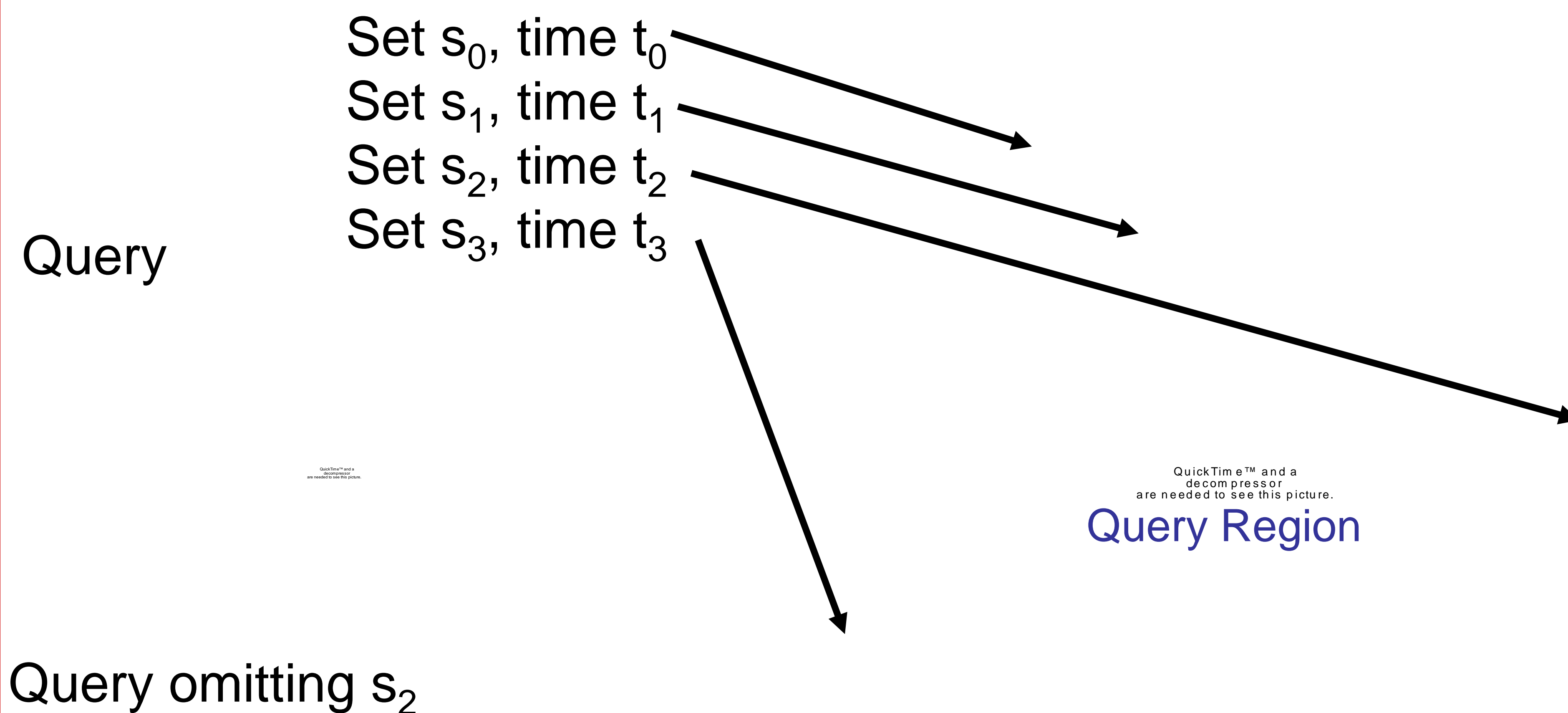


## Problem Definition and Motivation

- Sets of spatial data points added at distinct time epochs.
- Range searches only return most recent data where sets overlap.
- Irregular update regions may require alpha-shapes for definition.
- Queries can optionally omit subsets from consideration.
- Motivation: Queries on time-stamped massive data surveys.
- Challenge: Massive data requires I/O model; complex queries.



## Persistent Data Structures

- Persistent data structures maintain versioning history and data, retaining search complexity.
- Alterations tracked between versions of structure.
- Partially (only current version editable) or fully (all versions editable) persistent.
- Planar point location in  $O(\log n)$  query,  $O(n)$  space.
- Most persistent structures are non-spatial.

QuickTime™ and a decompressor are needed to see this picture.

## Planned Approaches

- Can simply treat time as an additional dimension.
- Requires complex query for stated results.
- Preferable: Make a persistent spatial data structure.
- Time treated as discrete epochs, making this approach more applicable.
- Can a 2-d data structure be made persistent to support rectangle + time epoch queries in  $O(\log N + K)$  time using  $O(N)$  space?
- Can updates be performed in  $O(M \log N)$  time for  $M$  points added at epoch  $t_i$ ?
- Can persistent spatial data structures be made I/O-efficient?

## Testing Plans

- Data sets: Shallow Survey 2008 Common Dataset (CCOM/JHC), as well as a synthetic data set.
- SSCD contains >580 GB of survey data; sets of binary point data range from 2 to 28 GB.
- Synthetic data set for initial testing randomly generated.

